Codewords of small weight in the (dual) code of points and k-spaces of $PG(n, q)$

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Abstract

In this paper, we study the p-ary linear code $C_k(n, q)$, $q = p^h$, p prime, $h \geq 1$, generated by the incidence matrix of points and k-dimensional spaces in $PG(n, q)$. We note that the condition $k \geq n/2$ arises in results on the codewords of $C_k(n, q)$. For $k \geq n/2$, we link codewords of $C_k(n, q) \setminus C_k(n, q)^\perp$ of weight smaller than $2q^k$ to k-blocking sets. We first prove that such a k-blocking set is uniquely reducible to a minimal k-blocking set, and exclude all codewords arising from small linear k-blocking sets. For $k < n/2$, we present counterexamples to lemmas valid for $k \geq n/2$. Next, we study the dual code of $C_k(n, q)$ and present a lower bound on the weight of the codewords, hence extending the results of Sachar [12] to general dimension.

1 Introduction

Let $PG(n, q)$ denote the n-dimensional projective space over the finite field $\mathbb{F}_q$ with q elements, where $q = p^h$, p prime, $h \geq 1$, and let $V(n+1, q)$ denote the underlying vector space. Let θ_n denote the number of points in $PG(n, q)$, i.e., $\theta_n = (q^{n+1} - 1)/(q - 1)$. A *blocking set* of $PG(n, q)$ is a set K of points such that each hyperplane of $PG(n, q)$ contains at least one point of K. A blocking set K is called *trivial* if it contains a line of $PG(n, q)$. These blocking sets are also called *1-blocking sets* in [3]. In general, a *k-blocking set* K in $PG(n, q)$ is a set of points such that any $(n - k)$-dimensional subspace intersects K. A k-blocking set K is called *trivial* when a k-dimensional subspace is contained in K. The smallest non-trivial k-blocking sets are characterized as cones with a $(k - 2)$-dimensional vertex π_{k-2} and a non-trivial 1-blocking set of minimum cardinality in a plane, skew to π_{k-2}, of $PG(n, q)$ as base curve [3, 8]. If an $(n - k)$-dimensional space contains exactly one point of a k-blocking set K in $PG(n, q)$, it is called a *tangent $(n - k)$-space* to K, and a point P of K is called *essential* when it belongs to a tangent $(n - k)$-space of K. A k-blocking set K is called *minimal* when no proper subset of K is also a k-blocking set, i.e., when each point of K is essential.

A lot of attention has been paid to blocking sets in the Desarguesian plane $PG(2, q)$, and to k-blocking sets in $PG(n, q)$. It follows from results of Sziklai

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[13], Szőnyi [14], and Szőnyi and Weiner [15] that every minimal k-blocking set K in $PG(n,q)$, $q = p^h$, p prime, $h \geq 1$, of size smaller than $3(q^{n-k}+1)/2$, intersects every subspace in zero or in 1 (mod p) points. If e is the largest integer such that K intersects every space in zero or 1 (mod p^e) points, then e is a divisor of h. This implies, for instance, that the cardinality of a minimal blocking set, of size smaller than $3(q+1)/2$, in $PG(2,q)$ can only lie in a number of intervals, each of which corresponds to a divisor e of h.

We define the incidence matrix $A = (a_{ij})$ of points and k-spaces in the projective space $PG(n,q)$, $q = p^h$, p prime, $h \geq 1$, as the matrix whose rows are indexed by the k-spaces of $PG(n,q)$ and whose columns are indexed by the points of $PG(n,q)$, and with entry

$$a_{ij} = \begin{cases} 1 & \text{if point } j \text{ belongs to } k\text{-space } i, \\ 0 & \text{otherwise.} \end{cases}$$

The p-ary linear code C of points and k-spaces of $PG(n,q)$, $q = p^h$, p prime, $h \geq 1$, is the $\mathbb{F}_p$-span of the rows of the incidence matrix A. From now on, we denote this code by C_k, or, if we want to specify the dimension and order of the ambient space, by $C_k(n,q)$. The *support* of a codeword c, denoted by $supp(c)$, is the set of all non-zero positions of c. The *weight* of c is the number of non-zero positions of c and is denoted by $wt(c)$. Often we identify the support of a codeword with the corresponding set of points of $PG(n,q)$. We let c_P denote the symbol of the codeword c in the coordinate position corresponding to the point P, and let (c_1,c_2) denote the scalar product in $\mathbb{F}_p$ of two codewords c_1, c_2 of C. Furthermore, if T is a subspace of $PG(n,q)$, then the incidence vector of this subspace is also denoted by T. The dual code $C^\perp$ is the set of all vectors orthogonal to all codewords of C, hence

$$C_k^\perp = \{v \in V(\theta_n, p) || (v,c) = 0, \ \forall c \in C_k\}.$$

This means that $(c,K) = 0$ for all k-spaces K of $PG(n,q)$. In [10], the p-ary linear code $C_{n-1}(n,q)$, $q = p^h$, p prime, $h \geq 1$, was discussed. The main goal of this paper is to prove similar results for the p-ary linear code $C_k(n,q)$ defined by the incidence matrix of points and k-spaces of $PG(n,q)$, $q = p^h$, p prime, $h \geq 1$. More precisely, in [10], the following results are proven.

Result 1. *(see also [1, Proposition 5.7.3]) The minimum weight codewords of $C_{n-1}(n,q)$ are the scalar multiples of the incidence vectors of the hyperplanes.*

Result 2. *There are no codewords with weight in the interval $]\theta_{n-1}, 2q^{n-1}[$ in $C_{n-1}(n,q)$, if q is prime, or if $q = p^2$, $p > 11$ prime.*

Result 3. *The only possible codewords of $C_{n-1}(n,q)$, with weight in the interval $]\theta_{n-1}, 2q^{n-1}[$, are the scalar multiples of non-linear minimal blocking sets.*

Result 4. *The minimum weight of $C_{n-1}(n,q) \cap C_{n-1}(n,q)^\perp$ is equal to $2q^{n-1}$.*

Result 5. *If c is a codeword of $C_{n-1}(n,q)^\perp$ of minimal weight, then $supp(c)$ is contained in a plane of $PG(n,q)$.*

Theorem 16(2) and Theorem 17 extend Result 1 and the first part of Result 2 to general dimension. However, the generalization of the second part of Result 2 in Theorem 18 and the generalization of Result 3 in Theorem 16(1) are weaker, due to the lack of a generalization of Result 4 in the case where q is not a prime. In Theorem 11, Result 5 is generalized.

In the study of codewords $c \in C_k(n,q)$ of weight smaller than $2q^k$, we distinguish the cases $c \in C_k(n,q) \setminus C_k(n,q)^\perp$ and $c \in C_k(n,q) \cap C_k(n,q)^\perp$. In the first case, for $k \geq n/2$, $supp(c)$ defines a k-blocking set of $PG(n,q)$. We eliminate the small linear k-blocking sets as possible codewords, if $k \geq n/2$. One of the results we need regarding k-blocking sets, is the unique reducibility property of k-blocking sets, of size smaller than $2q^k$, to a minimal k-blocking set. We derive this property in the next section.

2 A unique reducibility property for k-blocking sets in $PG(n,q)$ of size smaller than $2q^k$

In [14], algebraic curves are associated to blocking sets in $PG(2,q)$, in order to prove the following result.

Result 6. *[14, Szőnyi] If K is a blocking set in $PG(2,q)$ of cardinality $|K| \leq 2q$, then K can be reduced in a unique way to a minimal blocking set.*

In this section, we extend this result to general k-blocking sets in $PG(n,q)$, $n \geq 3$, by associating an algebraic hypersurface to a blocking set in $PG(n,q)$.

Let K be a blocking set in $PG(n,q)$, $n \geq 3$, with $|K| \leq 2q-1$. Suppose that the coordinates of the points are $(x_0,\ldots,x_n)$, where $X_n = 0$ defines the hyperplane at infinity H_∞, and let U be the set of affine points of K. Let $|K| = q+k+N$, $N \geq 1$, where N is the number of points of K in H_∞. Furthermore we assume that $(0,\ldots,0,1,0) \in K$. The hyperplanes not passing through $(0,\ldots,0,1,0)$ have equations $m_0X_0 + \cdots + m_{n-2}X_{n-2} - X_{n-1} + b = 0$ and they intersect H_∞ in the $(n-2)$-dimensional space $X_n = m_0X_0 + \cdots + m_{n-2}X_{n-2} - X_{n-1} = 0$. We call the $(n-1)$-tuple $\bar{m} = (m_0,\ldots,m_{n-2})$ the *slope* of the hyperplane. We also identify a slope $\bar{m}$ with the corresponding subspace $X_n = m_0X_0 + \cdots + m_{n-2}X_{n-2} - X_{n-1} = 0$ of dimension $n-2$ at infinity.

Definition 1. *Define the* Rédei polynomial *of U as*

$$\begin{aligned} H(X,X_0,\ldots,X_{n-2}) &= \prod_{(a_0,\ldots,a_{n-1})\in U} (X + a_0X_0 + \cdots + a_{n-2}X_{n-2} - a_{n-1}) \\ &= X^{q+k} + h_1(X_0,\ldots,X_{n-2})X^{q+k-1} + \cdots + \\ &\quad h_{q+k}(X_0,\ldots,X_{n-2}). \end{aligned}$$

For all $j = 1,\ldots,q+k$, $\deg h_j \leq j$. For simplicity of notations, we will also write $H(X,X_0,\ldots,X_{n-2})$ as $H(X,\bar{X})$.

Definition 2. *Let C be the affine hypersurface, of degree k, of $AG(n,q)$, defined by*

$$f(X,\bar{X}) = X^k + h_1(\bar{X}_i)X^{k-1} + \cdots + h_k(\bar{X}_i) = 0.$$

Theorem 1. *(1) For a fixed slope $\bar{m}$ defining an $(n-2)$-dimensional subspace at infinity not containing a point of K, the polynomial $X^q - X$ divides $H(X,\bar{m})$. Moreover, if $k < q-1$, then $H(X,\bar{m})/(X^q - X) = f(X,\bar{m})$ and $f(X,\bar{m})$ splits into linear factors over $\mathbb{F}_q$.*

(2) For a fixed slope $\bar{m} = (m_0, \ldots, m_{n-2})$, the element x is an r-fold root of $H(X, \bar{m})$ if and only if the hyperplane with equation $m_0X_0 + \cdots + m_{n-2}X_{n-2} - X_{n-1} + x = 0$ intersects U in exactly r points.

(3) If $k < q - 1$ and $\bar{m}$ defines an $(n-2)$-dimensional subspace at infinity not containing a point of K, such that the line $X_0 = m_0, \ldots, X_{n-2} = m_{n-2}$ intersects $f(X, \bar{X})$ at $(x, m_0, \ldots, m_{n-2})$ with multiplicity r, then the hyperplane with equation $m_0X_0 + \cdots + m_{n-2}X_{n-2} - X_{n-1} + x = 0$ intersects K in exactly $r + 1$ points.

Proof. (1) For every $X = b$, the hyperplane $m_0X_0 + \cdots + m_{n-2}X_{n-2} - X_{n-1} + b = 0$ contains a point $(a_0, \ldots, a_{n-1})$ of U. So $X - b$ is a factor of $H(X, \bar{m})$.

If $k < q - 1$, then $H(X, \bar{m}) = X^{q+k} + h_1(\bar{m})X^{q+k-1} + \cdots + h_{q+k}(\bar{m}) = (X^k + h_1(\bar{m})X^{k-1} + \cdots + h_k(\bar{m}))(X^q - X) = f(X, \bar{m})(X^q - X)$.

Since $H(X, \bar{m})$ splits into linear factors over $\mathbb{F}_q$, this is also true for $f(X, \bar{m})$.

(2) The multiplicity of a root $X = x$ is the number of linear factors in the product defining $H(X, \bar{m})$ that vanish at $(x, \bar{m})$. This is the number of points of U lying on the hyperplane $m_0X_0 + \cdots + m_{n-2}X_{n-2} - X_{n-1} + x = 0$.

(3) Since $(m_0, \ldots, m_{n-2})$ defines an $(n-2)$-dimensional subspace at infinity not containing a point of K, part (1) holds. If the intersection multiplicity is r, then x is an $(r+1)$-fold root of $H(X, \bar{m})$. Hence, the result follows from (2). □

Remark 1. *By induction on the dimension one can construct an $(n-2)$-dimensional subspace α skew to K. Since $|K| \leq 2q - 1$, K has a tangent hyperplane because all hyperplanes through α must contain at least one point of K.*

Assume that $X_n = 0$ is a tangent hyperplane to K in the point $(0, \ldots, 0, 1, 0)$. The following theorem links the problem of minimality of the blocking set K to that of the problem of finding linear factors of the affine hypersurface $C : f(X, \bar{X}) = 0$.

Theorem 2. *(1) If a point $P = (a_0, \ldots, a_{n-1}) \in U$ is not essential, then the linear factor $a_0X_0 + \cdots + a_{n-2}X_{n-2} - a_{n-1} + X$ divides $f(X, \bar{X})$.*

(2) If the linear factor $X + a_0X_0 + \cdots + a_{n-2}X_{n-2} - a_{n-1}$ divides $f(X, \bar{X})$, then $P = (a_0, \ldots, a_{n-1}) \in U$ and this point is not essential.

Proof. (1) Consider an arbitrary slope $\bar{m} = (m_0, \ldots, m_{n-2})$. For this slope $\bar{m}$, there are at least two points of K in the hyperplane $m_0X_0 + \cdots + m_{n-2}X_{n-2} - X_{n-1} + b = 0$ through $(a_0, \ldots, a_{n-1})$. Hence, by Theorem 1, the hyperplane $\pi : a_0X_0 + \cdots + a_{n-2}X_{n-2} - a_{n-1} + X = 0$ shares the point $(X, \bar{X}) = (a_{n-1} - (a_0m_0 + \cdots + a_{n-2}m_{n-2}), m_0, \ldots, m_{n-2})$ with C. Suppose that $a_0X_0 + \cdots + a_{n-2}X_{n-2} - a_{n-1} + X$ does not divide $f(X, \bar{X})$, and let R be a point of the hyperplane π not lying in C.

There are $q^{n-2} + \cdots + q + 1$ lines through R in the hyperplane π, and none of them is contained in C since $R \notin C$. Since such lines contain at most k points of C, π contains at most $k(q^{n-2} + \cdots + q + 1) < (q-1)(q^{n-2} + \cdots + q + 1) = q^{n-1} - 1$ points of C. This is a contradiction because of the number of possibilities for $\bar{m}$.

(2) If this linear factor divides $f(X, \bar{X})$, then for all $\bar{m} = (m_0, \ldots, m_{n-2})$, the hyperplane with slope $\bar{m}$ through $(a_0, \ldots, a_{n-1})$ intersects U in at least two points (Theorem 1 (3)). Here, we use that $X_n = 0$ is a tangent hyperplane to

K in the point $(0,\ldots,0,1,0)$, so $\bar{m}$ defines an $(n-2)$-dimensional subspace at infinity not containing a point of K.

Suppose that $(a_0,\ldots,a_{n-1}) \notin U$. By induction, it is possible to prove that there is a subspace π of dimension $n-2$ passing through $(a_0,\ldots,a_{n-1})$ and containing no points of K (cf. Remark 1). Consider all hyperplanes through π. One of them passes through $(0,\ldots,0,1,0)$; the other ones contain at least two points of K. So $|K| \geq 2q+1$, which is false.

Hence, $P = (a_0,\ldots,a_{n-1}) \in U$. Since all hyperplanes through P, including those through $(0,\ldots,0,1,0)$, contain at least two points of K, the point P is not essential. □

Corollary 1. *A blocking set B of size smaller than $2q$ in $PG(n,q)$ is uniquely reducible to a minimal blocking set.*

Proof. The non-essential points of B correspond to the linear factors over $\mathbb{F}_q$ of the polynomial $f(X,\bar{X})$, and this polynomial is uniquely reducible. □

We will extend this unique reducibility property to blocking sets with respect to k-blocking sets.

Theorem 3. *A k-blocking set in $PG(n,q)$ of size smaller than $2q^k$ is uniquely reducible to a minimal k-blocking set.*

Proof. Embed $PG(n,q)$ in $PG(n,q^k)$. Let π be a hyperplane of $PG(n,q^k)$. The space $\pi \cap \pi^q \cap \pi^{q^2} \cap \cdots \cap \pi^{q^{k-1}}$ is the intersection of π with $PG(n,q)$. Since it is the intersection of k hyperplanes, it has dimension at least $n-k$. This implies that a k-blocking set B in $PG(n,q)$ is also a 1-blocking set in $PG(n,q^k)$. In Corollary 1, it is proven that this latter blocking set is uniquely reducible to a minimal 1-blocking set B' in $PG(n,q^k)$. Since every $(n-k)$-dimensional space Π in $PG(n,q)$ can be extended to a hyperplane in $PG(n,q^k)$ that intersects $PG(n,q)$ only in Π (straightforward counting), it is easy to see that the minimal blocking set B' in $PG(n,q^k)$ is the unique minimal k-blocking set in $PG(n,q)$ contained in B. □

3 The linear code generated by the incidence matrix of points and k-spaces in $PG(n,q)$

In this section, we investigate the codewords of small weight in the p-ary linear code generated by the incidence matrix of points and k-dimensional spaces, or shortly *k-spaces*, in $PG(n,q)$, $q=p^h$, p prime, $h \geq 1$.

Lemma 1. *If U_1 and U_2 are subspaces of dimension at least $n-k$ in $PG(n,q)$, then $U_1 - U_2 \in C_k^{\perp}$.*

Proof. For every subspace U_i of dimension at least $n-k$ and every k-space K, $(K,U_i)=1$, hence $(K,U_1-U_2)=0$, so $U_1-U_2 \in C_k^{\perp}$. □

Note that in Lemma 1, $\dim U_1 \neq \dim U_2$ is allowed.

Lemma 2. *The scalar product (c,U), with $c \in C_k$ and U an arbitrary subspace of dimension at least $n-k$, is a constant.*

Proof. Lemma 1 yields that $U_1 - U_2 \in C_k^\perp$, for all subspaces U_1, U_2 with $\dim(U_i) \geq n-k$, hence $(c, U_1 - U_2) = 0$, so $(c, U_1) = (c, U_2)$. $\square$

Theorem 4. *The support of a codeword $c \in C_k$ with weight smaller than $2q^k$, for which $(c, S) \neq 0$ for some $(n-k)$-space S, is a minimal k-blocking set in $PG(n,q)$. Moreover, c is a codeword taking only values from $\{0, a\}$, $a \in \mathbb{F}_p^*$, and supp(c) intersects every $(n-k)$-dimensional space in $1 \pmod{p}$ points.*

Proof. If c is a codeword with weight smaller than $2q^k$, and $(c, S) = a \neq 0$ for some $(n-k)$-space, then, according to Lemma 2, $(c, S) = a$ for all $(n-k)$-spaces S, so *supp*(c) defines a k-blocking set B.

Suppose that every $(n-k)$-space contains at least two points of the k-blocking set B. Counting the number of incident pairs $(P \in B, (n-k)\text{-space through } P)$ yields that

$$|B| \begin{bmatrix} n \\ n-k \end{bmatrix} = \begin{bmatrix} n+1 \\ n-k+1 \end{bmatrix} 2.$$

Using that $|B| < 2q^k$ gives a contradiction. So there is a point $R \in B$ on a tangent $(n-k)$-space. Since c_R is equal to a, according to Lemma 2, $c_{R'} = a$ for every essential point R' of B.

Suppose B is not minimal, i.e. suppose there is a point $R \in B$ that is not essential. By induction on the dimension, we find an $(n-k-1)$-dimensional space π tangent to B in R. If every $(n-k)$-space through π contains two extra points of B, then $|B| > 2q^k$, a contradiction. Hence, there is an $(n-k)$-space S, containing besides R only one extra point R' of *supp*(c), such that $(c, S) = c_R + c_{R'} = a$. But since B is uniquely reducible to a minimal blocking set B (see Theorem 3), R' is essential, hence, $c_{R'} = a$. But this implies that $c_R = 0$, a contradiction. We conclude that the k-blocking set B is minimal.

Since all the elements R of *supp*(c) have the coordinate value $c_R = a$, and since $(c, H) = a$ for every $(n-k)$-dimensional space H, necessarily *supp*(c) intersects every $(n-k)$-dimensional space in $1 \pmod{p}$ points. $\square$

The following theorem is proved using the arguments from [15].

Theorem 5. *Let c be a codeword of $C_k(n,q)$ with weight smaller than $2q^k$, for which $(c, S) \neq 0$ for some $(n-k)$-space. Every subspace of $PG(n,q)$ that intersects supp(c) in at least one point, intersects it in 1 $\pmod{p}$ points.*

We emphasize that from now on, for some of the results, it is necessary to assume that $k \geq n/2$.

The following lemmas are extensions of the lemmas in [10]; we include the proofs to illustrate where the extra requirement $k \geq n/2$ arises.

Lemma 3. *(See [10, Lemma 3]) Assume $k \geq n/2$. A codeword c of C_k is in $C_k \cap C_k^\perp$ if and only if $(c, U) = 0$ for all subspaces U with $\dim(U) \geq n-k$.*

Proof. Let c be a codeword of $C_k \cap C_k^\perp$. Since $c \in C_k^\perp$, $(c, K) = 0$ for all k-spaces K, Lemma 2 yields that $(c, U) = 0$ for all subspaces U with dimension $n-k$ since $k \geq n-k$. Now suppose $c \in C_k$ and $(c, U) = 0$ for all subspaces U with dimension at least $n-k$. Applying this to a k-space yields that $c \in C_k \cap C_k^\perp$ since $k \geq n-k$. $\square$

Remark 2. *If $k < n/2$, the lemma is false. Let c be $K_1 - K_2$, with K_1 and K_2 two skew k-spaces. It is clear that $c \in C_k$ and that $(c, S) = 0$ for all $(n-k)$-spaces S. But $c \notin C_k^\perp$ since $(c, K_1) = 1 \neq 0$. Note that the lemma is still valid in one direction: if $c \in C_k \cap C_k^\perp$, then $(c, S) = 0$ for all $(n-k)$-spaces. For, let S be an $(n-k)$-space, and let K_i, $i = 1, \ldots, \theta_{n-2k}$, be the θ_{n-2k} k-spaces through a fixed $(k-1)$-space K' contained in S. Since $(c, K) = 0$ for all k-spaces K, it follows that $(c, S) = (c, K_1 \setminus K') + \cdots + (c, K_{\theta_{n-2k}} \setminus K') + (c, K') = 0$.*

Lemma 4. *(See [10, Lemma 4]) For $k \geq n/2$,*

$$C_k \cap C_k^\perp = \langle K_1 - K_2 || K_1, K_2 \text{ distinct } k\text{-spaces in } PG(n,q) \rangle .$$

Proof. Put $A = \{K_1 - K_2 || K_1, K_2 \text{ distinct } k\text{-spaces in } PG(n,q)\}$. Since $k \geq n/2$, two k-spaces K and K' of $PG(n,q)$ intersect in 1 (mod p) points, so $(K, K') = 1$. Hence, $A \subseteq C \cap C^\perp$, since $(K, v) = (K, K_i) - (K, K_j) = 1 - 1 = 0$, for every k-space K of $PG(n,q)$, and for every $v = K_i - K_j \in A$.

Moreover, since $\langle A \cup \{K_i\} \rangle$ contains each k-space, it follows that $\dim(C) - 1 \leq \dim(\langle A \rangle) \leq \dim(C \cap C^\perp)$. The lemma now follows easily, since $C \cap C^\perp$ is not equal to C, as a k-space, with $k \geq n/2$, is not orthogonal to itself. □

Remark 3. *If $k < n/2$, the theorem is false, since $K_1 - K_2 \notin C_k \cap C_k^\perp$, with K_1, K_2 two skew k-spaces (see Remark 2).*

The following lemmas are extensions of Lemmas 6.6.1 and 6.6.2 of Assmus and Key [1]. They will be used to exclude non-trivial small linear blocking sets as codewords. The proofs are an extension of the proofs of Lemmas 7 and 8 of [10].

Lemma 5. *For $k \geq n/2$, a vector v of $V(\theta_n, p)$ taking only values from $\{0, a\}$, $a \in \mathbb{F}_p^\star$, is contained in $(C_k \cap C_k^\perp)^\perp$ if and only if $|supp(v) \cap K| \pmod p$ is independent of the k-space K of $PG(n,q)$.*

Remark 4. *If $k < n/2$, the theorem is false. Let v be a k-space. It follows that $v \in (C_k \cap C_k^\perp)^\perp$ since $v \in C_k = (C_k^\perp)^\perp \subseteq (C_k \cap C_k^\perp)^\perp$. But $|supp(v) \cap K|$ is 0 $\pmod p$ or 1 $\pmod p$, depending on the k-space K.*

Lemma 6. *Assume $k \geq n/2$ and let c,v be two vectors taking only values from $\{0, a\}$, for some $a \in \mathbb{F}_p^\star$, with $c \in C_k$, $v \in (C_k \cap C_k^\perp)^\perp$. If $|supp(c) \cap K| \equiv |supp(v) \cap K| \pmod p$ for every k-space K, then $|supp(c) \cap supp(v)| \equiv |supp(c)| \pmod p$.*

As mentioned in the introduction, we will eliminate all so-called non-trivial *linear* k-blocking sets as the support of a codeword of C of small weight. In order to define a linear k-blocking set, we introduce the notion of a Desarguesian spread.

By what is sometimes called "field reduction", the points of $PG(n,q)$, $q = p^h$, p prime, $h \geq 1$, correspond to $(h-1)$-dimensional subspaces of $PG((n+1)h - 1, p)$, since a point of $PG(n,q)$ is a 1-dimensional vector space over $\mathbb{F}_q$, and so an h-dimensional vector space over $\mathbb{F}_p$. In this way, we obtain a partition $\mathcal{D}$ of the point set of $PG((n+1)h-1, p)$ by $(h-1)$-dimensional subspaces. In general, a partition of the point set of a projective space by subspaces of a given dimension k is called a *spread*, or a *k-spread* if we want to specify the dimension. The spread we have obtained here is called a *Desarguesian spread*. Note that the Desarguesian spread satisfies the property that each subspace spanned by two spread elements is again partitioned by spread elements. In fact, it can be shown that if $n \geq 2$, this property characterises a Desarguesian spread [11].

Definition 3. *Let U be a subset of $PG((n+1)h-1,p)$ and let $\mathcal{D}$ be a Desarguesian $(h-1)$-spread of $PG((n+1)h-1,p)$, then $\mathcal{B}(U) = \{R \in \mathcal{D} || U \cap R \neq \emptyset\}$.*

Analogously to the correspondence between the points of $PG(n,q)$ and the elements of a Desarguesian spread $\mathcal{D}$ in $PG((n+1)h-1,p)$, we obtain the correspondence between the lines of $PG(n,q)$ and the $(2h-1)$-dimensional subspaces of $PG((n+1)h-1,p)$ spanned by two elements of $\mathcal{D}$, and in general, we obtain the correspondence between the $(n-k)$-spaces of $PG(n,q)$ and the $((n-k+1)h-1)$-dimensional subspaces of $PG((n+1)h-1,p)$ spanned by $n-k+1$ elements of $\mathcal{D}$. With this in mind, it is clear that any hk-dimensional subspace U of $PG(h(n+1)-1,p)$ defines a k-blocking set $\mathcal{B}(U)$ in $PG(n,q)$. A blocking set constructed in this way is called a *linear k-blocking set*. Linear k-blocking sets were first introduced by Lunardon [11], although there a different approach is used. For more on the approach explained here, we refer to [9].

The following lemmas, theorems, and remarks are proven in the same way as the authors do in [10].

We put $N = hk$ throughout the following results. We call a linear k-blocking set B of $PG(n,q)$, $q = p^h$, p prime, $h \geq 1$, defined by an N-dimensional space of $PG(h(n+1)-1,p)$ a *small* linear k-blocking set.

Lemma 7. *Let U_N be an N-dimensional subspace of $PG(h(n+1)-1,p)$. The number of spread elements of $\mathcal{B}(U_N)$ intersecting U_N in exactly one point is at least $p^N - p^{N-2} - p^{N-3} - \cdots - p^{N-h+1} - p^{N-h-2} - \cdots - p^{N-2h+1} - p^{N-2h-2} - \cdots - p^{h+1} - p^{h-2} - \cdots - p$.*

Remark 5. *It follows from Lemma 7 that the number of spread elements of $\mathcal{B}(U_N)$ intersecting U_N in exactly one point is at least $p^N - p^{N-1} + 1$. We will use this weaker bound.*

Lemma 8. *If there are $p^N - p^{N-1} + 1$ points R_i of a minimal k-blocking set B in $PG(n,q)$, for which it holds that every line through R_i is either a tangent line to B or is entirely contained in B, then B is a k-space of $PG(n,q)$.*

Remark 6. *It follows from the proof of Lemma 8 in [10] that it is sufficient to find k linearly independent points R_i such that every line through R_i is either a tangent line to B or is entirely contained in B to prove that B is a k-space. Moreover, this bound is tight. If there are only $k-1$ linearly independent points for which this condition holds, we have the counterexample of a Baer cone, i.e. let B be the set of all lines connecting a point of a Baer subplane $\pi = PG(2,\sqrt{q})$ to the points of a $(k-2)$-dimensional subspace of $PG(n,q)$, skew to π.*

Theorem 6. *For every small linear k-blocking set B, not defining a k-space in $PG(n,p^h)$, there exists a small linear k-blocking set B' intersecting B in $2 \pmod{p}$ points.*

Using this, we exclude in Theorem 7 all small non-trivial linear k-blocking sets as codewords.

Theorem 7. *Assume $k \geq n/2$. If v is the incidence vector of a small non-trivial linear k-blocking set in $PG(n,q)$, then $v \notin C_k(n,q)$.*

Proof. We know that $|supp(v)| \equiv 1 \pmod p$ [10, Lemma 9]. We know from Theorem 6 that there exists a small linear k-blocking set w such that $|supp(v) \cap supp(w)| \equiv 2 \pmod p$. Since $|supp(w) \cap K| \equiv 1 \pmod p$ for every k-space K, see [10, Lemma 9], it follows that $w \in (C \cap C^{\perp})^{\perp}$ (Lemma 5). Similarly $|supp(v) \cap K| \equiv 1 \pmod p$, for every k-space K. Suppose that $v \in C$. Lemma 6 implies that $|supp(v) \cap supp(w)| \equiv |supp(v)| \pmod p \equiv 1 \pmod p$, a contradiction. □

Corollary 2. *For $k \geq n/2$, the only possible codewords c of $C_k(n,q)$ of weight in $]\theta_k, 2q^k[$, such that $(c,S) \neq 0$ for an $(n-k)$-space S, are scalar multiples of non-linear minimal k-blocking sets of $PG(n,q)$.*

Remark 7. *In view of Corollary 2 it is important to mention the conjectures made in [13]. If these conjectures are true (i.e. all small minimal blocking sets are linear), then Corollary 2 eliminates all codewords of $C_k(n,q) \backslash C_k(n,q)^{\perp}$ of weight in the interval $]\theta_k, 2q^k[$.*

For $q = p$ prime and for $q = p^2$, $p > 11$ prime, we can exclude all such possible codewords. We rely on the following results.

Theorem 8. *The only minimal k-blocking sets B in $PG(n,p)$, with p prime and $|B| < 2p^k$, such that every $(n-k)$-space intersects B in $1 \pmod p$ points, are k-spaces of $PG(n,p)$.*

Proof. By induction on the dimension, it is possible to prove that if a line contains at least two points of B, then this line is contained in B. It now follows, by induction on the dimension, that B is a k-space. □

To exclude codewords in $C_k(n,p^2)$, with p a prime, we can use the following theorem of Weiner which implies that every small minimal blocking set in $PG(n,p^2)$ is linear.

Theorem 9. *[16] A non-trivial minimal blocking set of $PG(n,p^2)$, $p > 11$, p prime, with respect to k-spaces and of size less than $3(p^{2(n-k)}+1)/2$ is a $(t, 2((n-k)-t-1))$-Baer cone with as vertex a t-space and as base a $2((n-k)-t-1)$-dimensional Baer subgeometry, where $\max\{-1, n-2k-1\} \leq t < n-k-1$.*

Theorems 8 and 9, together with Corollary 2, yield the following corollary.

Corollary 3. *There are no codewords c, with $wt(c) \in]\theta_k, 2q^k[$, in $C_k(n,q) \backslash C_k(n,q)^{\perp}$, with $k \geq n/2$, q prime or $q = p^2$, $p > 11$, p prime.*

4 The dual code of $C_k(n,q)$

In this section, we consider codewords c in the dual code $C_k(n,q)^{\perp}$ of $C_k(n,q)$. The goal of this section is to find a lower bound on the minimum weight of the code $C_k(n,q)^{\perp}$. Denote the minimum weight of a code C by $d(C)$.

In the following lemmas, the problem of finding the minimum weight of $C_k(n,q)^{\perp}$ is reduced to finding the minimum weight of $C_1(n-k+1,q)^{\perp}$. Note that $d(C_k(n,q)^{\perp}) \leq 2q^{n-k}$ since the difference of the incidence vectors of two $(n-k)$-spaces of $PG(n,q)$, intersecting in an $(n-k-1)$-space, is a codeword of $C_k(n,q)^{\perp}$.

Lemma 9. *For each $n \geq 2$, $0 < k \leq n-1$, the following inequalities hold:*

$$d(C_k(n,q)^{\perp}) \geq d(C_{k-1}(n-1,q)^{\perp}) \geq \cdots \geq d(C_1(n-k+1,q)^{\perp}).$$

Proof. Let c be a codeword of $C_k(n,q)^{\perp}$ of minimum weight, let R be a point of $PG(n,q)\backslash supp(c)$, lying in a tangent line to $supp(c)$, and let H be a hyperplane of $PG(n,q)$ not containing R. For each point $P \in H$, define $c'_P = \sum c_{P_i}$, with P_i the points of $supp(c)$ on the line $\langle R, P\rangle$, and let c' denote the vector with coordinates c'_P, $P \in H$. It easily follows that $c' \in C_{k-1}(n-1,q)^{\perp}$, and $supp(c')$ is contained in the projection of $supp(c)$ from the point R onto the hyperplane H. Clearly, $|supp(c')| \leq |supp(c)|$. Using this relation on a codeword c of minimum weight yields that $d(C_{k-1}(n-1,q)^{\perp}) \leq d(C_k(n,q)^{\perp})$. Continuing this process proves the statement. □

Theorem 10. *For each $n \geq 2$, $0 < k \leq n-1$, $d(C_k(n,q)^{\perp}) = d(C_1(n-k+1,q)^{\perp})$.*

Proof. Embed $\pi = PG(n-k+1,q)$ in $PG(n,q)$, $n > 2$, and extend each codeword c of $C_1(\pi)^{\perp}$ to a vector $c^{(n)}$ of $V(\theta_n,p)$ by putting a zero at each point $P \in PG(n,q)\backslash\pi$. Since the all one vector of $V(\theta_{n-k+1},p)$ is a codeword of $C_1(n-k+1,q)$, it follows that $\sum_{P\in\pi} c_P^{(n)} = 0$ for each $c^{(n)}$. This implies that $(c^{(n)},K) = 0$, for each k-space K of $PG(n,q)$ which contains π. If a k-space K of $PG(n,q)$ does not contain π, then $(c^{(n)}, K\cap\pi) = 0$, since $K \cap \pi$ is a line or can be described as a pencil of lines through a given point, and $(c,l) = 0$ for each line l of π. It follows that $c^{(n)}$ is a codeword of $C_k(n,q)^{\perp}$ of weight equal to the weight of c, which implies that $d(C_k(n,q)^{\perp}) \leq d(C_1(n-k+1,q)^{\perp})$. Regarding Lemma 9, this yields that $d(C_k(n,q)^{\perp}) = d(C_1(n-k+1,q)^{\perp})$. □

Lemma 10. *Let B be a set of points in $PG(n,q)$, with $\dim\langle B\rangle \geq n-k+2$, such that if a point in $PG(n,q)\backslash B$ lies on at least one secant line to B, then it does not lie on tangent lines to B, then $|B| \geq \theta_{n-k+1}$.*

Proof. We first prove the following result.

Let P be a point in B and let L be a line through P, lying in a plane π through P,R,S, with $R,S \in B$ and $P \notin RS$, then L is a secant line to B. If L is a tangent line to B, then the point $RS \cap L$ lies on a secant line and on a tangent line, a contradiction.

By induction, we prove that for each point $P \in B$, there exists an r-space π_r, with $r \leq n-k+2$, such that all lines through P in π_r are secant lines. The case $r = 2$ is already settled, so suppose that the statement is true for r, $r < n-k+2$. There is a point $T \in B \notin \pi_r$ since $\dim\langle B\rangle \geq n-k+2$. If M is a line through P in $\langle\pi_r, T\rangle$, then $\langle M, T\rangle$ intersects π_r in a line N through P, which is a secant line according to the induction hypothesis. Hence, we find three non-collinear points in B in the plane $\langle N, T\rangle$, so M is a secant line, so there is an $(r+1)$-space for which any line through P is a secant line. Counting the points of B on lines through P yields that $|B| \geq \theta_{n-k+1}$. □

Theorem 11. *If c is a codeword of $C_k(n,q)^{\perp}$, $n \geq 3$, of minimal weight, then $supp(c)$ is contained in an $(n-k+1)$-space of $PG(n,q)$.*

Proof. As already observed, we may assume that $wt(c) \leq 2q^{n-k}$. Assume that $\dim\langle supp(c)\rangle \geq n-k+2$. Using Lemma 10, we find a point $R \notin supp(c)$ lying on a tangent line to $supp(c)$ and lying on at least one secant line to $supp(c)$. It follows from Theorem 10 that

$$wt(c) = d(C_k(n,q)^\perp) = d(C_{k-1}(n-1,q)^\perp) = d(C_1(n-k+1,q)^\perp).$$

Let c' be defined as in the proof of Lemma 10. Since R lies on at least one secant line to $supp(c)$, $0 < wt(c') < wt(c)$. But this implies that c' is a codeword of $C_{k-1}(n-1,q)^\perp$ satisfying $0 < wt(c') \leq wt(c) - 1 < d(C_{k-1}(n-1,q)^\perp)$, a contradiction. □

In Theorem 11, we proved that finding the minimum weight of the code $C_k(n,q)^\perp$ is equivalent to finding the minimum weight of the code $C_1(n-k+1,q)^\perp$ of points and lines in $PG(n-k+1,q)$. Hence, we can use the following result due to Bagchi and Inamdar.

Result 7. *[2, Proposition 2] When q is prime, the minimum weight of the dual code $C_1(n,q)^\perp$ is $2q^{n-1}$. Moreover, the codewords of minimum weight are precisely the scalar multiples of the difference of two hyperplanes.*

Using Result 7, together with Theorem 11, yields the following theorem.

Theorem 12. *The minimum weight of $C_k(n,p)^\perp$, where p is a prime, is equal to $2p^{n-k}$, and the codewords of weight $2p^{n-k}$ are the scalar multiples of the difference of two $(n-k)$-spaces intersecting in an $(n-k-1)$-space.*

When q is not a prime, this result is false; we will present some counterexamples.

Theorem 13. *Let B be a minimal $(n-k)$-blocking set in $PG(n,q)$ of size $q^{n-k}+x$, with $x < (q^{n-k}+1)/2$, such that there exists an $(n-k)$-space T intersecting B in x points. The difference of the incidence vectors of B and T is a codeword of $C_k(n,q)^\perp$ with weight $2q^{n-k}+\theta_{n-k-1}-x$.*

Proof. If $x < (q^{n-k}+1)/2$, then B is a small minimal $(n-k)$-blocking set, hence every k-space intersects B in $1 \pmod p$ points (see [15]). Let c_1 be the incidence vector of B and let c_2 be the incidence vector of an $(n-k)$-space intersecting B in x points. Then $(c_1 - c_2, K) = (c_1, K) - (c_2, K) = 0$ for all k-spaces K, hence $c_1 - c_2$ is a codeword of $C_k(n,q)^\perp$, with weight $|B| + |T| - 2|B \cap T| = 2q^{n-k} + \theta_{n-k-1} - x$. □

We can use this theorem to lower the upper bound on the possible minimum weight of codewords of $C_k(n,q)^\perp$. Put $V(n+1,q) = V(1,q) \times V(n-k,q) \times V(k,q) = \mathbb{F}_q \times \mathbb{F}_{q^{n-k}} \times \mathbb{F}_{q^k}$ and put

$$B = \left\{(1, x, Tr(x)) || x \in \mathbb{F}_{q^{n-k}}\right\} \cup \left\{(0, x, Tr(x)) || x \in \mathbb{F}_{q^{n-k}}, x \neq 0\right\},$$

where Tr is the trace function of $\mathbb{F}_{q^{n-k}}$ to $\mathbb{F}_p$, p prime. The set B is a subset of $\mathbb{F}_q \times \mathbb{F}_{q^{n-k}} \times \mathbb{F}_q$ since $Tr(x) \in \mathbb{F}_p \subset \mathbb{F}_q, \forall x$. Moreover, B is a linear subspace, and therefore induces a Redéi type blocking set of size $q^{n-k} + (q^{n-k}-1)/(p-1)$, say $q^{n-k}+x$, w.r.t. the lines in $PG(\mathbb{F}_q \times \mathbb{F}_{q^{n-k}} \times \mathbb{F}_q) \cong PG(n-k+1,q)$, since there is an $(n-k)$-space π such that $|B \cap \pi| = x$. Embedding B in $PG(n,q)$ yields that B is a minimal blocking set w.r.t. k-spaces, hence B is a minimal $(n-k)$-blocking set such that there exists an $(n-k)$-space that intersects B in x points.

Using this, together with Theorem 13, yields the following corollary.

Corollary 4. *For $q = p^h$, p prime, $h \geq 1$,*

$$d(C_k(n,q)^\perp) \leq 2q^{n-k} + \theta_{n-k-1} - \frac{q^{n-k}-1}{p-1}.$$

In the case where q is even, [2] gives an upper bound on the minimum weight.

Result 8. *[2, Proposition 4] For q even, the minimum weight of the code $C_1(n,q)^\perp$ is at most $q^{n-2}(q+2)$.*

Result 8, together with Theorem 11, has the following corollary.

Corollary 5. *For q even, the minimum weight of $C_k(n,q)^\perp$ is at most $q^{n-k-1}(q+2)$.*

Remark 8. *It is easy to see that the minimum weight of $C_1(n-k+1,q)^\perp$, hence of $C_k(n,q)^\perp$, is at least $\theta_{n-k}+1$ since in $C_1(n-k+1,q)^\perp$, every line through a point of supp(c), with $c \in C_1(n-k+1,q)^\perp$, has to contain at least one other point of supp(c). If q is odd, Theorems 14 and 15 improve this lower bound. If q is even, then $d(C_k(n,q)^\perp) > \theta_{n-k}+1$, for $n > 3$, since otherwise, supp(c) would be a set B of points in $PG(n-k+1,q)$, no three collinear, and [7, Theorem 27.4.6] states that $|B| \leq q^{n-k} - q^{n-k-1}/2 + 4q^{n-7/2}$, a contradiction. For $n = 3$ and $k = 1$, [6, Lemma 16.1.4] yields that $|B| \leq q^2+1$, a contradiction. For $n = 3$ and $k = 2$, it is easy to see that the minimum weight is $q+2$.*

We will now prove a lower bound on the minimum weight of $C_k(n,q)^\perp$, q not a prime, q odd, by extending the bound of Sachar [12] on the minimum weight of $C_1(2,q)^\perp$.

Lemma 11. *Suppose that there are $2m$ different non-zero symbols used in the codeword $c \in C_k(n,q)^\perp$, q odd. Then*

$$wt(c) \geq \frac{4m}{2m+1}\theta_{n-k} + \frac{2m}{2m+1}.$$

Proof. We use the same techniques as in the proof of Proposition 2.2 in [12]. Let c be a codeword in $C_k^\perp$. Assume that $wt(c) \leq 2q^{n-k}$, and write $wt(c)$ as $\theta_{n-k} + x$.

Through every point P of *supp*(c), we can construct by induction on s, an s-space that only intersects *supp*(c) in P, through a fixed $(s-1)$-space only intersecting *supp*(c) in P, if $s \leq k-1$, since the number of s-spaces through an $(s-1)$-space is $(q^{n-s+1}-1)/(q-1) > 2q^{n-k}$ if $n-s > n-k$. So through every point P of *supp*(c), there is a $(k-1)$-space K' which intersects *supp*(c) only in the point P. For simplicity of notations, we use the terminology 2-*secant* for a k-space having two points of *supp*(c). Let $\bar{K}$ be a $(k-1)$-space intersecting *supp*(c) in one point, for which the number of 2-secants through $\bar{K}$ is minimal. We denote this number by X, or by X_R in case $\bar{K}$ intersects *supp*(c) in the point R of *supp*(c).

Since c is orthogonal to every k-space, if K is a 2-secant through R and R', R, $R' \in$ *supp*(c), then $c_R + c_{R'} = 0$, so the symbol $c_{R'}$ occurs at least X times in c. In fact, the number of occurrences of a certain non-zero symbol is always at least X.

The number of 2-secants through a given $(k-1)$-space intersecting *supp*(c) in exactly one point, is at least $\theta_{n-k} - x + 1$. So it is easy to see that the number of non-zero symbols used in c must be even; let this number of non-zero symbols be $2m$.

This implies that

$$2m(\theta_{n-k} - x + 1) \le \theta_{n-k} + x.$$

Hence,

$$x \ge \frac{2m-1}{2m+1}\theta_{n-k} + \frac{2m}{2m+1},$$

and

$$wt(c) \ge \frac{4m}{2m+1}\theta_{n-k} + \frac{2m}{2m+1}.$$

□

Theorem 14. *If $p \neq 2$, then $d(C_k(n,q)^\perp) \ge (4\theta_{n-k}+2)/3$, $q = p^h$, p prime, $h \ge 1$.*

Proof. Let c be a codeword of $C_k(n,q)^\perp$ with $wt(c) < (4\theta_{n-k}+2)/3$. According to Lemma 11, there is only one non-zero symbol used in c. Construct a $(k-1)$-space π through a point R of $supp(c)$ intersecting $supp(c)$ only in R. Then every k-space K through π has to contain at least $p-1$ extra points of $supp(c)$ in order to get $(c,K) = 0$. But then $wt(c) \ge (p-1)\theta_{n-k}+1$, a contradiction. □

Theorem 15. *The minimum weight of $C_k(n,q)^\perp$ is at least $(12\theta_{n-k}+2)/7$ if $p = 7$, and at least $(12\theta_{n-k}+6)/7$ if $p > 7$.*

Proof. We use the same techniques as in the proof of Proposition 2.4 in [12]. Let c be a codeword of minimum weight of $C_k(n,q)^\perp$ and suppose that $wt(c) < (12\theta_{n-k}+6)/7$. It follows from Lemma 11 that there are at most four different non-zero symbols used in the codeword c. Suppose first that there are exactly two non-zero symbols used in c, say 1 and -1. Suppose that the symbol -1 occurs the least, say y times. Construct a $(k-1)$-space π through a point R of $supp(c)$, where $c_R = 1$ and $\pi \cap supp(c) = \{R\}$. Every k-space $\bar{\pi}$ through π contains at least a second point of $supp(c)$. At most y of those k-spaces contain a point R' of $supp(c)$ with $c_{R'} = -1$, so at least $\theta_{n-k} - y$ of those k-spaces only contain points R' of $supp(c)$ with $c_{R'} = 1$. Since $(c, \bar{\pi}) = 0$, such k-spaces contain 0 (mod p) points of $supp(c)$. This yields

$$wt(c) \ge (\theta_{n-k} - y)(p-1) + y + 1.$$

Using that $wt(c) < (12\theta_{n-k}+2)/7$ implies that

$$p\theta_{n-k} - 7\theta_{n-k} - p + 7 < 0,$$

a contradiction if $p = 7$. Using that $wt(c) < (12\theta_{n-k}+6)/7$ implies that

$$(p-7)\theta_{n-k} + 7 - 3p < 0,$$

a contradiction if $p > 7$.

So we may assume that there are four non-zero symbols used in c, say $1, -1, a, -a$. Using the same notations as in the proof of Lemma 11, we see that

$$wt(c) \ge 4X_R. \tag{1}$$

We call a k-space through one of the $(k-1)$-spaces $\bar{K}$, with $\bar{K} \cap supp(c) = \{R\}$, that has exactly two extra points of $supp(c)$, a *3-secant*. Let X_3 denote the number of 3-secants through $\bar{K}$, and let X_w denote the number of k-spaces through $\bar{K}$ that intersect $supp(c)$ in more than 3 points. We have the following equations:

$$wt(c) \geq 1 + X_R + 2X_3 + 3X_w, \tag{2}$$
$$\theta_{n-k} = X_R + X_3 + X_w. \tag{3}$$

Suppose first that there are no 3-secants, then substituting (3) in (1) and (2) gives

$$wt(c) \geq 4\theta_{n-k} - 4X_w, \tag{4}$$
$$wt(c) \geq 1 + \theta_{n-k} + 2X_w. \tag{5}$$

Eliminating X_w using (4) and (5) gives

$$3wt(c) \geq 6\theta_{n-k} + 2,$$

a contradiction. This implies that $X_3 \neq 0$. Let T be a 3-secant through $\bar{K}$. The sum of the symbols used in T has to be zero, hence

$$(*) \begin{array}{l} 0 = 1 + 1 + a \text{ and } a = -2, \text{ or} \\ 0 = 1 + a + a \text{ and } a = -1/2. \end{array}$$

For each point P with $c_P = -a$, the k-space through $\bar{K}$ containing P has to intersect $supp(c)$ in more than three points, since otherwise

$$\begin{aligned} 1 - a - a &= 0 \text{ and } a = 1/2 \text{ or} \\ 1 + 1 - a &= 0 \text{ and } a = 2. \end{aligned}$$

This contradicts (*) since $p > 5$ implies that $\{2, -2\}$ cannot be the same as $\{1/2, -1/2\}$. There are at least X_R points with coefficient $-a$ and we see that they all must be on k-spaces contributing to X_w. Thus counting points again, we have

$$\begin{aligned} wt(c) &\geq 1 + X_R + 2X_3 + X_R \\ &= 1 + 2(\theta_{n-k} - X_3 - X_w) + 2X_3 \\ &= 1 + 2\theta_{n-k} - 2X_w. \end{aligned} \tag{6}$$

Substituting (3) in (1) and (2) gives

$$wt(c) \geq 4(\theta_{n-k} - X_3 - X_w) \tag{7}$$
$$wt(c) \geq 1 + \theta_{n-k} + X_3 + 2X_w. \tag{8}$$

Eliminating X_3 and X_w using (6), (7) and (8) yields

$$7wt(c) \geq 12\theta_{n-k} + 6$$

and the proof is complete. □

The second part of the following theorem is Corollary 5.7.5 of [1]. Here we give an alternative proof, similar to [2, Proposition 1].

Theorem 16. *(1) The only possible codewords of weight in* $]\theta_k, (12\theta_k+6)/7[$ *in* $C_k(n,q)$, $k \geq n/2$, $q = p^h$, $p > 7$ *prime,* $h \geq 1$, *are scalar multiples of incidence vectors of non-linear blocking sets.*

(2) The minimum weight of $C_k(n,q)$ *is* θ_k, *and a codeword of weight* θ_k *is a scalar multiple of the incidence vector of a* k*-space.*

Proof. (1) According to Lemma 2, there are two possibilities for a codeword $c \in C_k$ with $wt(c) < 2q^k$. Either $(c, S) \neq 0$ for every $(n-k)$-dimensional space S, and Corollary 2 yields that c is a scalar multiple of the incidence vector of a non-linear blocking set, or $(c, S) = 0$ for all $(n-k)$-spaces S. But this implies that $c \in C_{n-k}^{\perp}$, which has weight at least $(12\theta_k + 6)/7$ (see Theorem 15).

(2) For the second statement, it is sufficient to use a result of Bose and Burton [4] that shows that the minimum weight of a k-blocking set in $PG(n,q)$ is equal to θ_k, and that this minimum is reached if and only if the blocking set is a k-space. □

Remark 9. *In view of Theorem 16, it is important to mention the conjectures made in [13]. If these conjectures are true (i.e. all small minimal blocking sets are linear), then Theorem 16 eliminates all codewords of* $C_k(n,q)$ *of weight in the interval* $]\theta_k, (12\theta_k+6)/7[$.

In the cases $q = p$ and $q = p^2$, with p a prime, we can deduce more. Theorem 12, theorem 16, Theorem 8 and Theorem 9 yield the following theorems.

Theorem 17. *There are no codewords with weight in* $]\theta_k, 2q^k[$ *in* $C_k(n,q)$, $k \geq n/2$, *where* $q = p$ *is prime.*

Theorem 18. *There are no codewords with weight in* $]\theta_k, (12\theta_k+6)/7[$ *in* $C_k(n,q)$, $k \geq n/2$, *where* $q = p^2$, $p > 11$ *prime.*

We now turn our attention to codewords in $C_k(n,q)$, $k \geq n/2$, $q = p^h$, p prime, $h \geq 3$, with weight in $]\theta_k, (12\theta_k+6)/7[$. We know from Theorem 15 that such codewords belong to $C_k(n,q) \setminus C_k(n,q)^{\perp}$, so they define minimal k-blocking sets B intersecting every $(n-k)$-dimensional space in $1 \pmod{p}$ points (see Theorem 4, Lemma 3). Let e be the maximal integer for which B intersects every $(n-k)$-space in $1 \pmod{p^e}$ points. In [5, Corollary 5.2], it is proven that

$$|B| \geq q^k + \frac{q^k}{p^e+1} - 1.$$

We now derive an upper bound on $|B|$, based on [5, Theorem 5.3].

Theorem 19. *Let* B *be a minimal* k*-blocking set in* $PG(n,q)$, $n \geq 2$, $q = p^h$, p *prime,* $h \geq 1$, *intersecting every* $(n-k)$*-dimensional space in* $1 \pmod{p^e}$ *points, with* e *the maximal integer for which this is true. If* $|B| \in]\theta_k, (12\theta_k+6)/7[$ *and that* $p^e > 2$, *then*

$$|B| \leq q^k + \frac{2q^k}{p^e}.$$

Proof. Put $E = p^e$ and let τ_{1+iE} be the number of $(n-k)$-dimensional spaces intersecting B in $1 + iE$ points. We count the number of $(n-k)$-dimensional spaces, the number of incident pairs (R, π), with $R \in B$ and with π an $(n-k)$-dimensional space through R, and the number of triples (R, R', π), with R and

R' distinct points of B and π an $(n-k)$-dimensional space passing through R and R'. This gives us the following formulas.

$$\sum_{i\geq 0}\tau_{1+iE} = \frac{(q^{n+1}-1)(q^n-1)}{(q^{n-k+1}-1)(q^{n-k}-1)}\cdot X, \qquad (9)$$

$$\sum_{i\geq 0}(1+iE)\tau_{1+iE} = |B|\left(\frac{q^n-1}{q^{n-k}-1}\right)\cdot X, \qquad (10)$$

$$\sum_{i\geq 0}(1+iE)(1+iE-1)\tau_{1+iE} = |B|(|B|-1)\cdot X, \qquad (11)$$

where

$$X = \frac{(q^{n-1}-1)\cdots(q^{k+1}-1)}{(q^{n-k-1}-1)\cdots(q-1)}$$

is the number of $(n-k)$-dimensional spaces through a line of $PG(n,q)$. Since $\sum_{i\geq 0} i(i-1)E^2\tau_{1+iE}\geq 0$, we obtain

$$|B|(|B|-1)-(1+E)|B|\cdot\left(\frac{q^n-1}{q^{n-k}-1}\right)+(1+E)\left(\frac{(q^{n+1}-1)(q^n-1)}{(q^{n-k+1}-1)(q^{n-k}-1)}\right)\geq 0.$$

Under the condition $2<E$, this implies that

$$|B|\leq q^k+\frac{2q^k}{E}.$$

□

Remark 10. *If $p^e > 4$, then $|B| < 3/2q^k$ in which case results of Sziklai prove that e is a divisor of h [13, Corollary 5.2].*

We summarize the results on the minimum weight of $C_k(n,q)^\perp$, $k\geq n/2$, in the following table (with $\theta_n = (q^{n+1}-1)/(q-1)$).

p	h	d
2	$(k,n)\neq(n-1,n)$	$\theta_{n-k}+1<d\leq q^{n-k-1}(q+2)$
p	1	$2p^{n-k}$
$2<p<7$	$h>1$	$(4\theta_{n-k}+2)/3\leq d\leq 2q^{n-k}+\theta_{n-k-1}-\frac{q^{n-k}-1}{p-1}$
7	$h>1$	$(12\theta_{n-k}+2)/7\leq d\leq 2q^{n-k}+\theta_{n-k-1}-\frac{q^{n-k}-1}{p-1}$
$p>7$	$h>1$	$(12\theta_{n-k}+6)/7\leq d\leq 2q^{n-k}+\theta_{n-k-1}-\frac{q^{n-k}-1}{p-1}$

Table 1: The minimum weight d of $C_k(n,q)^\perp$, $k\geq n/2$, $q=p^h$, p prime, $h\geq 1$

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