# A geometric approach to determine an optimal 2-dimensional flow on a graph

Joint work with Davide Mattiolo, Jozef Rajník and Gloria Tabarelli

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• *G* – Simple graph



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 $(\mathbb{Z}_6,+)$ 

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- Kirchhoff's law satisfied at each vertex
- $\Gamma$ -flow  $(O, \varphi)$  on G



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Example of a nowhere-zero 4-flow



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#### 5-Flow Conjecture (Tutte, 1954)

Every bridgeless graph admits a nowhere-zero 5-flow.

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#### 5-Flow Conjecture (Tutte, 1954)

Every bridgeless graph admits a nowhere-zero 5-flow.

#### Remark

Constant 5 cannot be improved both for  $\mathbb{Z}\mbox{-flows}$  and for  $\mathbb{R}\mbox{-flows}.$ 

A *d*-dimensional nowhere-zero *r*-flow ((r, d)-NZF) is an  $\mathbb{R}^d$ -flow using only vectors with Euclidean norm from [1, r - 1]

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- d = 2: complex flows [Thomassen, 2014]

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- *d* = 1: circular nowhere-zero *r*-flows
- d = 2: complex flows [Thomassen, 2014]
- r = 2: Unit vector flows [Wang et. al, Thomassen]



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Is  $1 + \sqrt{2}$  the best possible value for  $K_4$ ?

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Is  $1 + \sqrt{2}$  the best possible value for  $K_4$ ? YES, but we have not a trivial proof

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The *d*-dimensional flow number of a bridgeless graph G is

 $\phi_d(G) = \inf\{r \in \mathbb{R} \mid G \text{ has } (r, d) \text{-NZF}\}$ 

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- continuous function  $x \mapsto$  maximum norm

(2,3)-NZF of the Petersen graph P ( $\phi_3(P) = 2$ ).

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(2,3)-NZF of the Petersen graph  $P(\phi_3(P) = 2)$ .



•  $z_1, z_2, z_3, z_4, z_5, x_1, x_2$  are unit vectors in  $\mathbb{R}^3$  $(z_1 + z_2 + z_3 + z_4 + z_5 = 0)$ 

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- $x_2 + z_3, x_2 + z_3 + z_4, x_2 + z_3 + z_4 + z_5, x_2 + z_3 + z_4 + z_5 + z_1$ are unit vectors in  $\mathbb{R}^3$

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For every bridgeless graph G:
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- $\phi_d(G)$  nonincreasing in d
- $\phi_1(G) \le 6$  [Seymour, 1981]

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- $\phi_1(G) \leq 6$  [Seymour, 1981]
- $\phi_7(G) = 2$  (observed by Thomassen [2014])

Relation to other conjectures:

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## Conjecture (Berge-Fulkerson 1971)

Each bridgeless cubic graph G has 6 perfect matchings such that every edge of G is contained in exactly two of them. Or, equivalently, G has a 6-cycle 4-cover.

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*BF* Conjecture  $\Rightarrow \phi_6(G) = 2$
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### Conjecture (5-Cycle Double Cover, Seymour/Szekeres 1979)

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Each bridgeless graph has 5-cycle 2-cover.

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#### Conjecture (5-Cycle Double Cover, Seymour/Szekeres 1979)

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Each bridgeless graph has 5-cycle 2-cover.

$$5 - CDC$$
 Conjecture  $\Rightarrow \phi_5(G) = 2$ 

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•  $\phi_1(G) \leq 5$  [Tutte, 1954]

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•  $\phi_1(G) \le 5$  [Tutte, 1954] •  $\phi_d(G) = 2, \forall d \ge 3$  [Jain, 2007] •  $\phi_2(G) < ???$ 

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- $\phi_1(G) \le 5$  [Tutte, 1954]
- $\phi_d(G) = 2, \forall d \ge 3$  [Jain, 2007]
- $\phi_2(G) \leq ???$  This is our starting point!

# Multi 2-dimensional flows

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### Theorem [Thomassen, 2014]

For a graph G, the following statements are equivalent

- (i) G has a (3,1)-NZF,
- (ii) G has a (2,2)-NZF,
- (iii) G has a (2,2)-NZF with values from  $\{z \in \mathbb{C} \mid z^3 = 1\}$ .

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Moreover, if G is cubic, then all of them are also equivalent to (iv) G is bipartite.

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Moreover, if G is cubic, then all of them are also equivalent to (iv) G is bipartite.

For this kind of questions an answer for cubic graphs implies an answer for the general case! Proposition (D.Mattiolo, G.M., J.Rajnik, G.Tabarelli, 2021)

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Theorem (Mattiolo, M., Rajnik, Tabarelli, 2021)

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Proof:

• Seymour 6-flow theorem  $(\phi_1(G) \leq 6)$ 

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Proof:

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- G has a 2-flow  $\varphi_2$  and a 3-flow  $\varphi_3$  (zeros allowed)
- $\forall e \in E(G) \colon \varphi_2(e) \neq 0 \text{ or } \varphi_3(e) \neq 0$

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- $\varphi(e) = (\varphi_2(e), \varphi_3(e))$
- $1 \le ||\varphi(e)|| \le \sqrt{1^2 + 2^2} = \sqrt{5}$

### Theorem (Mattiolo, M., Rajnik, Tabarelli)

If G has an oriented k-cycle double cover, then

• 
$$\phi_2(G) = 2$$
, if  $k = 3$ ;

• 
$$\phi_2(G) \le 1 + \sqrt{2}$$
, if  $k = 4$ ;

• 
$$\phi_2(G) \leq \Phi^2$$
, if  $k = 5$ . (where  $\Phi = \frac{1+\sqrt{5}}{2}$ )

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- The proof is based on finding a set of k points in the complex plane such that the ratio of maximum distance to minimum distance is the smallest possible.
- For k = 3, 4, 5 the best possible configurations are proved to be the vertices of a regular k-gon. [Bateman, Erdös (1951)]

### Oriented 5-CDC Conjecture (Jaeger, 1988)

Each bridgeless graph has a collection of five oriented cycles such that each edge is contained in exactly two of them, once in each direction.

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#### Corollary (Mattiolo, M., Rajnik, Tabarelli, 2021)

If the oriented 5-CDC conjecture holds true, then  $\phi_2(G) \leq \Phi^2 \approx 2.618$  for every bridgeless graph G.

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An oriented 5-CDC of the Petersen graph



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#### A 5-NZF of the Petersen graph

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#### A 5-NZF of the Petersen graph

It is well-known that  $\phi_1(P) = 5$ 

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A ( $\Phi^2$ , 2)-NZF of the Petersen graph (Recall  $1 + \Phi = \Phi^2$ )

Looking for a graph with large 2-dimensional flow number: a better 2-dimensional flow on the Petersen graph

We use a different geometric construction to improve previous result.

Best upper bound known so far

$$\phi_2(P) \le 1 + \sqrt{7/3} \approx 2.5275$$

$$(\Phi^2 \approx 2.6180)$$

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## Flow triangulations

• Triangle: set of points (sides and interior)



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  - parallel
  - same length
  - triangles on opposite sides



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An *r*-flow triangulation of a graph G is a collection of triangles  $T_v$ ,  $\forall v \in V(G)$ :

• each edge incident with v corresponds to a unique side of  $T_v$ 



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- side lengths from [1, r 1]



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#### Examples of flow triangulations



Flow triangulations of  $K_4$  and  $K_{3,3}$ 

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#### Proposition (Mattiolo, M., Rajnik, Tabarelli)

A bridgeless cubic graph G has an r-flow triangulation if and only if G has an (r, 2)-flow.

# $\phi_2(P) \le 1 + \sqrt{7/3} \approx 2.5275$



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Is a nice *r*-flow triangulation possible for every (r, 2)-flow of *G*?

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• What if G is bipartite?

## Once again: General upper bound for $\phi_2(G)$ ?

• New possible conjecture:  $\phi_2(G) \leq 1 + \sqrt{7/3} \approx 2.5275$  for all G?

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- Computer assisted:  $\phi_2(P_{\Delta}) \leq 2.590296429$

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- It seems that φ<sub>2</sub>(P<sub>Δ</sub>) > φ<sub>2</sub>(P)
- Computer assisted:  $\phi_2(P_\Delta) \leq 2.590296429$
- We conjecture:  $\phi_2(G) \leq \Phi^2 \approx 2.6180$

#### 1-dimensional flows and 3-edge-colourings

#### 1-dimensional flow number of cubic graphs

If G is a cubic graph, then on of the following holds:

- $\phi_1(G) = 3$  and G is bipartite
- $\phi_1(G) = 4$  and G is 3-edge-colourable (non-bipartite)

•  $\phi_1(G) > 4$  and G is not 3-edge-colourable

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#### 2-dimensional flows and 3-edge-colourings

No such classification for  $\phi_2$ :

•  $\phi_2(J_5) \le 2.387893647 < 1 + \sqrt{2} = \phi_2(K_4)$ 



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#### Theorem - Mattiolo, M., Rajnik, Tabarelli 2022

Let G be a cubic graph and let g be its odd-girth. Then,

$$\phi_2(G) \geq \begin{cases} 1+2\sin(\frac{\pi}{6}\cdot\frac{g}{g-1}) & \text{if } g \equiv 1,3 \mod 6, \\ 1+2\sin(\frac{\pi}{6}\cdot\frac{g+1}{g}) & \text{if } g \equiv 5 \mod 6. \end{cases}$$

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## A general lower bound for $\phi_2(G)$ ?

#### $\phi_2(G) \geq \phi_2(W_g)$





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#### Theorem - Mattiolo, M., Rajnik, Tabarelli 2022

Let  $W_g$  be a wheel of order g + 1. Then,

$$\phi_2(W_g) = \begin{cases} 2 & \text{if } g \text{ even,} \\ 1 + 2\sin(\frac{\pi}{6} \cdot \frac{g}{g-1}) & \text{if } g \equiv 1,3 \mod 6, \\ 1 + 2\sin(\frac{\pi}{6} \cdot \frac{g+1}{g}) & \text{if } g \equiv 5 \mod 6. \end{cases}$$

Proof (more than 15 pages): the following three types of configurations are optimal



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## **Open Problems**

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• 
$$\phi_2(P) = 1 + \sqrt{7/3}?$$



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•  $\phi_2(G) \leq \Phi^2$  for every bridgeless graph G?

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- $\phi_2(G) \leq \Phi^2$  for every bridgeless graph G?
- Does a graph G such that  $\phi_2(G) = \Phi^2$  exist?

# Thanks for your attention

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